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A FINITE CHARACTERIZATION AND RECOGNITION OF INTERSECTION GRAPHS OF HYPERGRAPHS WITH RANK AT MOST 3 AND MULTIPLICITY AT MOST 2 IN THE CLASS OF THRESHOLD GRAPHS

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Abstract

We characterize the class L_3^2 of intersection graphs of hypergraphs with rank at most 3 and multiplicity at most 2 by means of a finite list of forbidden induced subgraphs in the class of threshold graphs. We also give an O(n)time algorithm for the recognition of graphs from L_3^2 in the class of threshold graphs, where n is the number of vertices of a tested graph.

Keywords: intersection graph, hypergraph rank, hypergraph multiplicity, forbidden induced subgraph, threshold graph.

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1. INTRODUCTION

In this paper, we consider finite undirected graphs without loops and multiple edges. The vertex and the edge sets of a graph G are denoted by V(G) and E(G), respectively; $N(v) = N_G(v)$ is the neighborhood of a vertex v in G and $\deg(v)$ is the degree of v; the subgraph of G induced by a set $X \subseteq V(G)$ is denoted by G(X). A vertex v of a graph G is called *dominating* if $N(v) \cup \{v\} = V(G)$.

The intersection graph $L(\mathcal{H})$ of a hypergraph \mathcal{H} is defined as follows:

- (1) the vertices of $L(\mathcal{H})$ are in a bijective correspondence with the edges of \mathcal{H} ;
- (2) two vertices are adjacent in $L(\mathcal{H})$ if and only if the corresponding edges have a non-empty intersection.

The rank of a hypergraph \mathcal{H} is the maximum size of its edges. The *multiplicity* of a pair of vertices u, v of \mathcal{H} is the number of edges in \mathcal{H} containing both u and v; the *multiplicity* $m(\mathcal{H})$ of \mathcal{H} is the maximum multiplicity among all pairs of vertices in \mathcal{H} (see for example [15]).

Denote by L_r^m the class of intersection graphs of hypergraphs with rank at most r and multiplicity at most m. So, we refer to L_r^∞ as the class of intersection graphs of hypergraphs with rank at most r. The class L_r^m , where $r \ge 1$, $m \ge 1$ or $m = \infty$, is hereditary (i.e., every induced subgraph of a graph in L_r^m is also in L_r^m). Therefore, it can be characterized by means of a list (finite or not) of forbidden induced subgraphs.

A non-trivial characterization of the class L_r^m is known only for $r \leq 2$. These are:

- Beineke's finite characterization of the class L_2^1 of line graphs (i.e., intersection graphs of simple graphs) [1];
- a finite characterization of the class L_2^{∞} of intersection graphs of multigraphs by Bermond and Meyer [2];
- a finite characterization of the class L_2^m by Tashkinov [22].

Such finite characterizations of the classes above imply that there exist polynomial algorithms for recognizing graphs from these classes. (For efficient algorithms for recognizing graphs from L_2^1 see, e.g., [4, 11, 17, 19].) It is also known that for any $r \geq 3$ and m, where $m \geq 1$ or $m = \infty$, there does not exist a finite characterization for the class L_r^m (see [6, 15, 16, 10]).

Poljak, Rödl and Turzik [18] proved that the problem of determining whether a graph belongs to L_r^{∞} is NP-complete for an arbitrary r. Moreover, they proved that for every fixed $r \geq 4$, the analogous problem remains NP-complete. The question whether or not the class L_3^{∞} can be recognized in polynomial time is still open, but recognizing intersection graphs of hypergraphs without multiple edges with rank at most 3 is NP-complete as well [18]. The following result generalizing one from [18] was obtained in [7]: For every fixed $m \geq 1$ and an arbitrary r, the problem of determining whether a graph belongs to L_r^m is NP-complete. Hliněný and Kratochvíl [8] proved that for every fixed $r \geq 3$, the problem of determining whether a graph belongs to L_r^1 is NP-complete. The class L_3^1 was studied in different papers, and several graph classes were found, where the problem of recognizing graphs from the class is polynomially solvable or remains NP-complete ([7, 9, 14, 15, 16, 21]).

A graph G is called *split* [5] if there exists a partition of its vertex set $V(G) = A \cup B$ into a clique A and a stable set B (*bipartition* (A, B)). It was proved in [12] that for every fixed r, there exists a finite characterization of the graphs from L_r^1 in the class of split graphs. In [13] (see also [7]), this result was generalized to the class L_r^m for every fixed m.

A split graph with the bipartition (A, B) is called *threshold* [3] if the vertices from B can be numbered as b_1, b_2, \ldots, b_k so that $N(b_1) \supseteq N(b_2) \supseteq \cdots \supseteq N(b_k)$. In [20], the problem of determining the Krausz dimension of a graph (the minimum r such that the graph belongs to the class L_r^1) was solved in the subclass of threshold graphs of the form $K_n - E(K_p)$.

In Section 2 of this paper, we give some preliminary facts (e.g., a so-called Krausz type characterization of the class L_3^2 in terms of clique coverings), prove some technical lemmas and formulate Theorem 2 that gives a finite characterization of the class L_3^2 (consisting of 15 graphs) in the class of threshold graphs. In Sections 3 and 4, we prove the necessity and sufficiency of Theorem 2, respectively. In Section 5 we give an O(n)-time algorithm for the recognition of graphs from L_3^2 in the class of threshold graphs, where n is the number of vertices of a tested graph.

2. Some Preliminaries and the Formulation of Theorem 2

A finite family $\mathscr{C} = (C_1, C_2, \ldots, C_q)$ of cliques of the graph G is called a *covering* of G if every vertex as well as every edge of G is contained in some C_i . The cliques C_i are the *clusters* of \mathscr{C} . For a vertex $v \in V(G)$, denote by $\mathscr{C}(v)$ the subfamily of all clusters of \mathscr{C} that contain v. A covering \mathscr{C} of the graph G is called an (r, m)-covering if any vertex of G belongs to at most r clusters of \mathscr{C} , and any two clusters of \mathscr{C} have at most m vertices in common.

Theorem 1 [7, 13]. A graph G belongs to the class L_3^2 if and only if there exists a (3, 2)-covering of G.

A clique of a graph G is called *maximal* if it is not contained in some other clique of G.

Let a threshold graph with the bipartition (A, B) be given, where $B = \{b_1, b_2, \ldots, b_k\}$ and $N(b_1) \supseteq N(b_2) \supseteq \cdots \supseteq N(b_k)$. We denote such a graph by $G(p, q_1, q_2, \ldots, q_k)$ if |A| = p and $\deg(b_i) = q_i$ for any $i = 1, 2, \ldots, k$. Without loss of generality (W.l.o.g.), we assume below that any threshold graph

 $G(p, q_1, q_2, \ldots, q_k)$ with the bipartition (A, B) satisfies the conditions $A = \{a_1, a_2, \ldots, a_p\}$, $B = \{b_1, b_2, \ldots, b_k\}$, $p > q_1$ and $N(b_i) = \{a_1, a_2, \ldots, a_{q_i}\}$ for any $i = 1, 2, \ldots, k$ (see Figure 1).



Figure 1. The graph G(3, 2, 1) and its bipartition (A, B).

In this paper, we characterize the class L_3^2 by means of a finite list of forbidden induced subgraphs in the class of threshold graphs:

Theorem 2. A threshold graph H belongs to the class L_3^2 if and only if it contains none of the graphs $K_{1,4}$, G(12,7), G(11,10), G(10,9,5), G(10,9,7), G(10,9,9), G(10,7,k), k = 1, 2, ..., 7, G(9,8,1), G(9,8,2) as induced subgraphs.

Now we formulate some technical statements that will be used for proving Theorem 2.

A (3,2)-covering $\mathscr{C} = (C_1, C_2, \ldots, C_t)$ of a complete graph G is called a decomposition (3,2)-covering if $C_i \neq V(G)$ for any $i = 1, 2, \ldots, t$.

Lemma 3. Let $\mathscr{C} = (C_1, C_2, \ldots, C_t)$ be a decomposition (3, 2)-covering of a complete graph G. Then the following statements hold:

- (i) $|C_i| \le 6$ for any i = 1, 2, ..., t.
- (ii) If $C_i \setminus C_j \neq \emptyset$ for some $i, j \in \{1, 2, \dots, t\}$, then $|C_j \setminus C_i| \leq 4$.
- (iii) If $(C_i \cap C_j) \setminus C_k \neq \emptyset$ for some different $i, j, k \in \{1, 2, \dots, t\}$, then $|C_k \setminus (C_i \cup C_j)| \leq 2$.

Proof. (i) Let, to the contrary, $C_i = \{a_1, a_2, \ldots, a_7, \ldots\}$ for some $i \in \{1, 2, \ldots, t\}$. Consider a vertex $v \in V(G) \setminus C_i$. By the definition of a (3, 2)-covering, each cluster of \mathscr{C} contains at most two edges of va_s , $s = 1, 2, \ldots, 7$. Hence, the edges va_s , $s = 1, 2, \ldots, 7$, are covered by at least four clusters of \mathscr{C} , and, therefore, the vertex v is contained in at least four clusters of \mathscr{C} , which is a contradiction to the definition of \mathscr{C} . (ii) Assume, to the contrary, that for a vertex $v \in V(G)$, we have $v \in C_i \setminus C_j$ and $C_j \setminus C_i = \{a_1, a_2, a_3, a_4, a_5, \ldots\}$. By the definition of a (3, 2)-covering, the edges va_s , $s = 1, 2, \ldots, 5$, are covered by at least three clusters of \mathscr{C} , different from C_i . So, taking into account the cluster C_i , the vertex v is contained in at least four clusters of \mathscr{C} , which is a contradiction to the definition of \mathscr{C} .

(iii) Let, instead, $v \in (C_i \cap C_j) \setminus C_k \neq \emptyset$ and $C_k \setminus (C_i \cup C_j) = \{a_1, a_2, a_3, \ldots\}$. By the definition of a (3,2)-covering, the edges va_1, va_2, va_3 are covered by at least two clusters of \mathscr{C} , different from C_i and C_j . So, together with the clusters C_i , C_j , the vertex v is contained in at least four clusters of \mathscr{C} , which is a contradiction.

Lemma 4. Let $\mathscr{C} = (C_1, C_2, \ldots, C_t)$ be a decomposition (3, 2)-covering of a complete graph G. Then the following statements hold:

- (i) If G has order 11, then it contains no cluster of size at most 2.
- (ii) If G has order 12, then it contains no cluster of size at most 3.

Proof. (i) Let $V(G) = \{a_1, a_2, \ldots, a_{11}\}, C_1 \in \mathscr{C}(a_1) \text{ and } |C_1| \leq 2$. W.l.o.g., assume that $\{a_3, a_4, \ldots, a_{11}\} \subseteq V(G) \setminus C_1$. By the definition of \mathscr{C} , there exists a cluster $C_2 \in \mathscr{C}(a_1)$ of size at least 6 among the clusters covering some of the nine edges a_1a_i , $i = 3, 4, \ldots, 11$. By Lemma 3(i),(ii), $|C_2| = 6$ and $C_1 \subseteq C_2$. Hence, $|V(G) \setminus (C_1 \cup C_2)| = 5$ and there exists a cluster $C_3 \in \mathscr{C}(a_1) \setminus \{C_1, C_2\}$ of size at least 6 containing the set $V(G) \setminus (C_1 \cup C_2)$. By Lemma 3(i), $C_3 = \{a_1\} \cup (V(G) \setminus (C_1 \cup C_2))$. We have $|C_2| = |C_3| = 6$ and $|C_2 \cap C_3| = 1$, which is a contradiction to Lemma 3(ii).

The statement (ii) of the lemma follows immediately from the statement (i).

3. Proof of Necessity of Theorem 2

By heredity of the class L_3^2 , one has to show that none of the graphs $K_{1,4}$, G(12,7), G(11,10), G(10,9,5), G(10,9,7), G(10,9,9), G(10,7,k), $k = 1, 2, \ldots, 7$, G(9,8,1) and G(9,8,2) belongs to this class. Obviously, there exists no (3,2)-covering for the star $K_{1,4}$. Therefore, $K_{1,4} \notin L_3^2$ by Theorem 1.

Furthermore, let G be one of the graphs G(12,7), G(11,10), G(10,9,5), G(10,9,7), G(10,9,9), G(10,7,k), k = 1, 2, ..., 7, G(9,8,1), G(9,8,2) with the bipartition (A, B). Suppose, to the contrary, that there exists a (3,2)-covering $\mathscr{D} = (D_1, D_2, ..., D_t)$ of G.

W.l.o.g., we will assume that no cluster of \mathscr{D} is contained in some other cluster of \mathscr{D} . By Theorem 1, it can be easily seen that $D_i \neq A$ for any $i = 1, 2, \ldots, t$, since $\deg(b_1) \geq 7$.

Put $\mathscr{C} = (C_1, C_2, \ldots, C_t)$, where $C_i = D_i \cap A$, $i = 1, 2, \ldots, t$. Then \mathscr{C} is a decomposition (3, 2)-covering of the subgraph G(A), since $N(b_i) \neq A$ for each

 $b_i \in B$. A cluster $C \in \mathscr{C}$ is called b_i -reduced with $b_i \in B$, if $C \cup \{b_i\} \in \mathscr{D}$. A cluster $C \in \mathscr{C}$ is called simply reduced if it is b_i -reduced for some $b_i \in B$. By Lemma 3(i), \mathscr{C} contains two or three b_1 -reduced clusters, since $\deg(b_1) \geq 7$.

Lemma 5. The following statements hold:

- (i) If $C_1, C_2 \in \mathscr{C}$ are two different b_i -reduced clusters with $b_i \in B$, then $|C_1 \cap C_2| \leq 1$.
- (ii) If $C_1, C_2 \in \mathscr{C}$ are two different b_i -reduced clusters with $b_i \in B$, then $C_1 \nsubseteq C_2$ and $C_2 \nsubseteq C_1$.
- (iii) If $C_1, C_2, C_3 \in \mathscr{C}$ are three different reduced clusters, then $C_1 \cap C_2 \cap C_3 = \emptyset$.

Proof. (i) The validity of the statement follows immediately from the definition of \mathscr{C} .

(ii) The statement follows from the above assumption that no cluster of \mathscr{D} is contained in some other cluster of \mathscr{D} .

(iii) If, to the contrary, $a \in C_1 \cap C_2 \cap C_3$, then the edge aa_p is not covered by a cluster from $\mathscr{C}(a) = \{C_1, C_2, C_3\}$, which is a contradiction to the definition of \mathscr{C} .

We consider the following separate cases and come to a contradiction in each of them.

(1) G = G(12, 7).

(a) Assume that there exist exactly two b_1 -reduced clusters $C_1, C_2 \in \mathscr{C}$. By Lemma 4(ii), $|C_1| \ge 4$ and $|C_2| \ge 4$. Hence, by Lemma 5(i) and the equality $|C_1 \cup C_2| = 7$, we obtain $|C_1| = |C_2| = 4$ and $|C_1 \cap C_2| = 1$. W.l.o.g., assume that $C_1 \cap C_2 = \{a_1\}$. Consider the cluster $C_3 \in \mathscr{C}(a_1) \setminus \{C_1, C_2\}$. Then $\{a_1, a_8, a_9, a_{10}, a_{11}, a_{12}\} \subseteq C_3$. By Lemma 3(i), $C_3 = \{a_1, a_8, a_9, a_{10}, a_{11}, a_{12}\}$ (see Figure 2). We have $|C_3 \setminus C_1| = 5$, which is a contradiction to Lemma 3(ii).



Figure 2. The clusters C_1 , C_2 and C_3 of the covering \mathscr{C} in the case (1).

(b) Suppose that there exist exactly three b_1 -reduced clusters $C_1, C_2, C_3 \in \mathscr{C}$. Taking into account Lemmas 5(i) and 4(ii), we obtain that $|C_1 \cup C_2| \ge 7$ and, therefore, $|C_1 \cup C_2 \cup C_3| \ge 9 > 7 = \deg(b_1)$, which is a contradiction.

(2) G = G(11, 10).

(a) Assume that there exist exactly two b_1 -reduced clusters $C_1, C_2 \in \mathscr{C}$. By Lemma 5(i), $|C_1 \cap C_2| \leq 1$. By Lemmas 5(ii) and 3(ii), $|C_1 \setminus C_2| \leq 4$ and $|C_2 \setminus C_1| \leq 4$. Therefore, $\deg(b_1) = |C_1 \cup C_2| \leq 9$, which is a contradiction.

(b) Let \mathscr{C} contain three b_1 -reduced clusters C_1 , C_2 and C_3 .

First, we suppose that C_1 , C_2 and C_3 are pairwise disjoint. By Lemmas 3(ii) and 4(i), we have $3 \leq |C_i| \leq 4$ for any i = 1, 2, 3. W.l.o.g., assume that $C_1 = \{a_1, a_2, a_3\}, C_2 = \{a_4, a_5, a_6\}, C_3 = \{a_7, a_8, a_9, a_{10}\}$. By the definition of \mathscr{C} and Lemma 3(i), we have $|\mathscr{C}(a_1)| = 3$, since $|A \setminus C_1| = 8$.

Let C_4 and C_5 be two clusters in $\mathscr{C}(a_1) \setminus \{C_1\}$. Each of the clusters C_4 and C_5 has at least one common vertex with any of the clusters C_2, C_3 . If, for example, $C_4 \cap C_2 = \emptyset$, then $a_1 \in (C_1 \cap C_4) \setminus C_2$ and $|C_2 \setminus (C_1 \cup C_4)| = |C_2| = 3$, which is a contradiction to Lemma 3(iii). Since $C_3 \subseteq C_4 \cup C_5$ by the definition of \mathscr{C} and $|C_3| = 4$, then each of the clusters C_4 and C_5 has exactly two common vertices with the cluster C_3 .

The inequalities $|C_4| \ge 5$ and $|C_5| \ge 5$ hold. Otherwise, let, for example, $|C_4| \le 4$. Then $|C_5| \ge 6$, since $|C_4 \cup C_5| \ge 9$. Hence, by Lemma 3(i), $|C_5| = 6$. Therefore, $C_4 \cap C_5 = \{a_1\}$ and $|C_5 \setminus C_4| = 5$, which is a contradiction to Lemma 3(ii).

W.l.o.g., assume that $\{a_4, a_7, a_8\} \subseteq C_4$, $\{a_6, a_9, a_{10}, a_{11}\} \subseteq C_5$. Since $|C_5 \setminus C_1| \leq 4$ by Lemma 3(ii), then $a_5 \notin C_5$. Hence, $a_5 \in C_4$. We have $a_5 \in (C_2 \cap C_4) \setminus C_5$. By Lemma 3(ii), $|C_5 \setminus (C_2 \cup C_4)| \leq 2$. Then $a_{11} \in C_4$ and, by Lemma 3(i), $C_4 = \{a_1, a_4, a_5, a_7, a_8, a_{11}\}$ (see Figure 3). Therefore, $|C_4 \setminus C_1| = 5$, which is a contradiction to Lemma 3(ii).

Now, w.l.o.g., assume that $a_1 \in C_1 \cap C_2$. By Lemma 5(i), $C_1 \cap C_2 = \{a_1\}$. By Lemmas 5(ii) and 3(ii), $|C_1| \leq 5$ and $|C_2| \leq 5$. Each of the clusters C_1 , C_2 has size at least 4. If not, then $a_1 \in (C_1 \cap C_2) \setminus C_3$ by Lemma 5(iii), and $|C_3 \setminus (C_1 \cup C_2)| \geq 10 - (3 + 5 - 1) = 3$, which is a contradiction to Lemma 3(iii).

Furthermore, assume that at least one of the clusters C_1 , C_2 , say C_1 , has size 5. Then $|C_1 \setminus C_3| \leq 4$ by Lemmas 5(ii) and 3(ii), and so $|C_1 \cap C_3| = 1$ by Lemma 5(i). Let $C_1 \cap C_3 = \{a_2\}$. Then $a_2 \in (C_1 \cap C_3) \setminus C_2$ by Lemma 5(ii). We obtain that $|C_2 \setminus (C_1 \cup C_3)| \leq 2$ by Lemma 3(iii). Therefore, $|C_2 \cap C_3| = 1$. Let $C_2 \cap C_3 = \{a_3\}$. We have $a_3 \in (C_2 \cap C_3) \setminus C_1$ and $|C_1 \setminus (C_2 \cup C_3)| = 3$, contradicting Lemma 3(iii).

Thus, $|C_1| = |C_2| = 4$. Let, w.l.o.g., $C_1 = \{a_1, a_2, a_3, a_4\}$, $C_2 = \{a_1, a_5, a_6, a_7\}$. Then $\{a_8, a_9, a_{10}\} \subseteq C_3$, since $\{a_1, a_2, \ldots, a_{10}\} = N(b_1)$. By Lemma 5(iii), $a_1 \in (C_1 \cap C_2) \setminus C_3$. However, then $|C_3 \setminus (C_1 \cup C_2)| = 3$, which is a contradiction to Lemma 3(iii).



Figure 3. The clusters C_1 , C_2 , C_3 , C_4 and C_5 of the covering \mathscr{C} in the case (2).

(3) G = G(10, 9, 5).

Each vertex a_i , where i = 1, 2, ..., 5, belongs to one b_1 - and one b_2 -reduced clusters. Therefore, by Lemma 5(iii), each two of the b_2 -reduced clusters have no common vertices. By Lemma 5(iii), if a vertex belongs to two of the b_1 -reduced clusters, then this vertex belongs to the set $\{a_6, a_7, a_8, a_9\}$.

(a) Let \mathscr{C} contain exactly two b_1 -reduced clusters C_1, C_2 . Since $|C_1 \cup C_2| = 9$, we get $|C_1 \cap C_2| = 1$ and $|C_1| = |C_2| = 5$ by Lemmas 5(i),(ii) and 3(ii). Let, w.l.o.g., $C_1 \cap C_2 = \{a_9\}$. By the definition of \mathscr{C} , any vertex a_i , where $i = 1, 2, \ldots, 8$, belongs to exactly two clusters from $\mathscr{C}(a_i) \setminus \{C_1, C_2\}$. Moreover, it is easy to obtain that, for any vertex a_i , where $i = 1, 2, \ldots, 8$, each cluster $C \in$ $\mathscr{C}(a_i) \setminus \{C_1, C_2\}$ satisfies the equalities $|C \cap (C_1 \setminus C_2)| = 2$ and $|C \cap (C_2 \setminus C_1)| = 2$. Since every b_2 -reduced cluster is a subset of $(C_1 \setminus C_2) \cup (C_2 \setminus C_1)$ and belongs to $\mathscr{C}(a_i) \setminus \{C_1, C_2\}$, it has size 4, which is a contradiction.

(b) Let \mathscr{C} contain three pairwise non-intersecting b_1 -reduced clusters C_1 , C_2 and C_3 . By Lemma 3(ii), $|C_i| \leq 4$ for every i = 1, 2, 3.

(b1) First, suppose that $|C_1| = 1$, $|C_2| = 4$ and $|C_3| = 4$. Put $C_1 = \{a_1\}$. Consider the clusters $C_4, C_5 \in \mathscr{C}(a_1) \setminus \{C_1\}$. By the definition of $\mathscr{C}, |C_i \cap C_j| = 2$ for any i = 2, 3 and j = 4, 5. In particular, $(C_4 \cap C_5) \cap (C_2 \cup C_3) = \emptyset$. Since $(C_2 \cap C_4) \setminus C_5 \neq \emptyset$, then $|C_5 \setminus (C_2 \cup C_4)| \leq 2$ by Lemma 3(iii). Similarly, $|C_4 \setminus (C_2 \cup C_5)| \leq 2$. Therefore, $a_{10} \in C_4 \cap C_5$. We obtain that there does not exist a b_2 -reduced cluster in $\mathscr{C}(a_1)$, which is a contradiction.

Now, let $C_1 \subset \{a_6, a_7, a_8, a_9\}$. W.l.o.g., put $C_1 = \{a_9\}$. Note that each b_2 -reduced cluster C in \mathscr{C} has size at most 4. If not (i.e., $|C| = \deg(b_2) = 5$), then the inclusion $C \subseteq C_2 \cup C_3$ implies that $|C \cap C_2| \ge 3$ or $|C \cap C_3| \ge 3$, which is

a contradiction to the definition of \mathscr{C} . Let C_4 be a b_2 -reduced cluster in \mathscr{C} with size at most 2. Let $a_1 \in C_4 \cap C_2$. Consider the cluster $C_5 \in \mathscr{C}(a_1) \setminus \{C_2, C_4\}$. By the definition of \mathscr{C} , we have $C_3 \setminus C_4 \subseteq C_5$. Since $|C_4| \leq 2$ and $C_4 \cap C_2 \neq \emptyset$, we have $|C_3 \setminus C_4| \geq 3$. Therefore, $|C_3 \cap C_5| \geq 3$, which is a contradiction.

(b2) Suppose that $|C_1| = 2$, $|C_2| = 3$ and $|C_3| = 4$. Let $a \in C_1$, where $a \in \{a_1, a_2, \ldots, a_9\}$. Consider the clusters $C_4, C_5 \in \mathscr{C}(a) \setminus \{C_1\}$. By the definition of \mathscr{C} , $1 \leq |C_i \cap C_2| \leq 2$ and $|C_i \cap C_3| = 2$ for any i = 4, 5. Moreover, at least one of the clusters C_4, C_5 , say C_5 , has exactly two common vertices with C_2 . Clearly, $(C_4 \cap C_5) \cap C_3 = \emptyset$ and $|(C_4 \cap C_5) \cap C_2| \leq 1$. If $a_{10} \in C_5$, then $|C_5| = 6$ by Lemma 3(i). We have $C_1 \setminus C_5 \neq \emptyset$ and $|C_5 \setminus C_1| = 5 > 4$, which is a contradiction to Lemma 3(ii). Therefore, $a_{10} \in C_4 \setminus C_5$. By Lemma 3(i), at least one vertex a' of the set $C_5 \cap C_2$ does not belong to C_4 . We obtain that $a' \in (C_2 \cap C_5) \setminus C_4$ and $|C_4 \setminus (C_2 \cup C_5)| \geq 3$, which is a contradiction to Lemma 3(ii).

(b3) Let $|C_1| = |C_2| = |C_3| = 3$. Assume that there exists a b_2 -reduced cluster in \mathscr{C} with size at most 2. Therefore, this cluster does not intersect with some of the clusters C_1 , C_2 and C_3 , which is a contradiction to the definition of \mathscr{C} .

Now, let $C_4 = N(b_2)$ be the only b_2 -reduced cluster in \mathscr{C} . W.l.o.g., assume that $C_1 = \{a_1, a_6, a_7\}, C_2 = \{a_2, a_3, a_8\}$ and $C_3 = \{a_4, a_5, a_9\}$. Consider the clusters $C' \in \mathscr{C}(a_2) \setminus \{C_2, C_4\}$ and $C'' \in \mathscr{C}(a_3) \setminus \{C_2, C_4\}$. By the definition of \mathscr{C} , we have $a_6, a_7, a_9, a_{10} \in C' \cap C''$. Therefore, C' = C''. Put $C_5 = C'$. Then $C_3 \setminus C_5 \neq \emptyset$ and $|C_5 \setminus C_3| = 5 > 4$, which is a contradiction to Lemma 3(ii).

(c) Let \mathscr{C} contain three b_1 -reduced clusters C_1, C_2, C_3 and $C_1 \cap C_2 \neq \emptyset$. W.l.o.g., assume that $|C_1| \geq |C_2|$. By Lemma 5(iii), we obtain that $(C_1 \cap C_2) \setminus C_3 \neq \emptyset$. It follows from Lemma 3(iii) that $|C_3 \setminus (C_1 \cup C_2)| \leq 2$. Hence, $|C_1 \cup C_2| \geq 7$. Then $|C_1| \geq 4$. Moreover, by Lemmas 5(ii) and 3(ii), we have $|C_1| \leq 5$.

(c1) Let $|C_1| = 5$. Then $C_1 \cap C_3 \neq \emptyset$ by Lemmas 5(ii) and 3(ii). Furthermore, $C_2 \cap C_3 = \emptyset$ by Lemmas 5(iii) and 3(iii). Since $(C_1 \cap C_3) \setminus C_2 \neq \emptyset$ and, by Lemma 3(iii), $|C_2 \setminus (C_1 \cup C_3)| \le 2$, we have $|C_1| = 5$, $|C_2| = 3$ and $|C_3| = 3$. Recall that $C_1 \cap C_2, C_1 \cap C_3 \subseteq \{a_6, a_7, a_8, a_9\}$. W.l.o.g., assume that $C_1 \cap C_2 =$ $\{a_8\}, C_1 \cap C_3 = \{a_9\}$. Consider the clusters $C_4 \in \mathscr{C}(a_8) \setminus \{C_1, C_2\}$ and $C_5 \in \mathscr{C}(a_8) \setminus \{C_1, C_2\}$ $\mathscr{C}(a_9) \setminus \{C_1, C_3\}$. By the definition of \mathscr{C} , we have $|C_4 \cap (C_1 \setminus \{a_8, a_9\})| \leq 1$ and $|C_5 \cap (C_1 \setminus \{a_8, a_9\})| \leq 1$. Note that $C_4 \cap (C_2 \setminus \{a_8\}) = \emptyset$. If, to the contrary, $a \in C_4 \cap (C_2 \setminus \{a_8\})$, then $\mathscr{C}(a) = \{C_2, C_4, C_5\}$ and some vertex of the set $C_1 \setminus \{a_8, a_9\}$ does not belong to the set $C_2 \cup C_4 \cup C_5$, contradicting the definition of \mathscr{C} . Analogously, $C_5 \cap (C_3 \setminus \{a_9\}) = \emptyset$. At least one of the clusters C_2, C_3 , say C_3 , contains a vertex $a' \in \{a_1, a_2, ..., a_5\}$, since $|\{a_1, a_2, ..., a_5\} \cap C_1| \leq 3$. Let a'' be another vertex in the set $C_3 \setminus \{a_9\}$. Consider the clusters $C' \in \mathscr{C}(a') \setminus \{C_3, C_4\}$ and $C'' \in \mathscr{C}(a'') \setminus \{C_3, C_4\}$. Each of them contains the set $(C_2 \setminus \{a_8\}) \cup (C_1 \setminus (C_3 \cup C_4))$ of size at least 4. Therefore, $C' = C'' = C_6$ is a cluster of \mathscr{C} of size at least 6. By Lemma 3(i), $|C_6| = 6$ (see Figure 4). Since $a' \in \{a_1, a_2, ..., a_5\}$, then C_6 is a b_2 -reduced cluster in \mathscr{C} , which is a contradiction.



Figure 4. The clusters C_1 , C_2 , C_3 , C_4 , C_5 and C_6 of the covering \mathscr{C} in the case (3).

(c2) Now, let $|C_1| = 4$. Then, taking into consideration the inequalities $|C_1 \cup C_2| \ge 7$ and $|C_1| \ge |C_2|$, we have $|C_2| = 4$.

Let C_3 intersect with C_1 or C_2 . Then, by Lemma 3(iii), C_3 intersects with both C_1 and C_2 . By Lemma 5(i),(iii), we can assume, w.l.o.g., that $C_1 =$ $\{a_1, a_2, a_7, a_8\}, C_2 = \{a_3, a_4, a_7, a_9\}$ and $C_3 = \{a_5, a_6, a_8, a_9\}$. Consider the cluster $C_4 \in \mathscr{C}(a_7) \setminus \{C_1, C_2\}$. By the definition of \mathscr{C} , we have $a_5, a_6, a_{10} \in C_4$. Initially, let $C_4 = \{a_5, a_6, a_7, a_{10}\}$. Consider the cluster $C_5 \in \mathscr{C}(a_5) \setminus \{C_3, C_4\}$. By the definition of \mathscr{C} , we have $a_1, a_2, a_3, a_4 \in C_5$. Both clusters C_3, C_4 are not b_2 reduced since each of them contains at least one of the vertices $a_6, a_7, a_8, a_9, a_{10}$. Hence C_5 is a b_2 -reduced cluster. It follows from the inclusion $N(b_2) \subseteq C_5$ that $C_5 = N(b_2) = \{a_1, a_2, \dots, a_5\}$. Consider the cluster $C_6 \in \mathscr{C}(a_6) \setminus \{C_3, C_4\}$. By the definition of \mathscr{C} , we have $a_1, a_2, a_3, a_4 \in C_6$. Thus $C_6 \neq C_5$ and $|C_6 \cap C_5| \geq C_5$ 4 > 2, which is a contradiction to the definition of \mathscr{C} . If the cluster C_4 has a non-empty intersection with the set $(C_1 \setminus C_2) \cup (C_2 \setminus C_1)$, for example $a_1 \in C_4$, then at least one of the vertices a_3, a_4 also belongs to C_4 . Otherwise, by the definition of \mathscr{C} , the cluster $C_5 \in \mathscr{C}(a_1) \setminus \{C_1, C_4\}$ contains the vertices a_3, a_4 and a_9 . We obtain that $C_5 \neq C_2$ and $|C_5 \cap C_2| \geq 3 > 2$, which is a contradiction. Let $a_3 \in C_4$ and $C_5 \in \mathscr{C}(a_1) \setminus \{C_1, C_4\}$. Then $a_4, a_9 \in C_5$. We obtain that none of the clusters $C_1, C_4, C_5 \in \mathscr{C}(a_1)$ is b_2 -reduced, which is a contradiction.

Assume that the cluster C_3 does not intersect with C_1 and C_2 . Then $|C_3| = 2$. One of the vertices a_6, a_7, a_8 and a_9 , say a_9 , belongs to $C_1 \cap C_2$. Consider the cluster $C_4 \in \mathscr{C}(a_9) \setminus \{C_1, C_2\}$. Clearly, $C_3 \cup \{a_{10}\} \subseteq C_4$. We show that $|C_4 \cap$ $(C_1 \setminus C_2)| = 1$ and $|C_4 \cap (C_2 \setminus C_1)| = 1$. Indeed, if C_4 has no common vertices with one of the sets $C_1 \setminus C_2$ or $C_2 \setminus C_1$, say with $C_1 \setminus C_2$, then $(C_3 \cap C_4) \setminus C_1 \neq \emptyset$ and $|C_1 \setminus (C_3 \cup C_4)| = 3$, contradicting Lemma 3(iii). Let $C_3 = \{a', a''\}$. Consider the clusters $C' \in \mathscr{C}(a') \setminus \{C_3, C_4\}$ and $C'' \in \mathscr{C}(a'') \setminus \{C_3, C_4\}$. We have $(C_1 \setminus C_4) \cup (C_2 \setminus C_4) \subseteq C' \cap C''$. Since $|(C_1 \setminus C_4) \cup (C_2 \setminus C_4)| = 4$, we obtain that C' = C'' by the definition of \mathscr{C} . Denote the cluster C' by C_5 . It can be easily obtained by the definition of \mathscr{C} that there are two clusters C_6 and C_7 in \mathscr{C} such that $((C_1 \setminus C_2) \cap C_4) \cup (C_2 \cap C_5) \cup \{a_{10}\} \subseteq C_6$ and $((C_2 \setminus C_1) \cap C_4) \cup (C_1 \cap C_5) \cup \{a_{10}\} \subseteq C_7$. Each vertex from the set $(C_1 \setminus C_2) \cup (C_2 \setminus C_1)$ belongs to exactly three of the non- b_2 -reduced clusters $C_1, C_2, C_4, C_5, C_6, C_7$. Clearly, at least three of the vertices a_1, a_2, \ldots, a_5 belong to the set $(C_1 \setminus C_2) \cup (C_2 \setminus C_1)$, which is a contradiction.

(4) We can come to a contradiction for each of the graphs G = G(10, 9, 9)and G = G(10, 9, 7) analogously to the graph G = G(10, 9, 5).

(5) $G = G(10, 7, k), k = 1, 2, \dots, 7.$

(a) First, assume that $4 \leq k \leq 7$. For any i = 1, 2, 3, 4, denote by C_{i1} and C_{i2} , respectively, b_1 - and b_2 -reduced clusters from $\mathscr{C}(a_i)$. Consider the cluster $C_{i3} \in \mathscr{C}(a_i) \setminus \{C_{i1}, C_{i2}\}$. Since $C_{i1}, C_{i2} \subseteq \{a_1, a_2, \ldots, a_7\}$, we have $\{a_8, a_9, a_{10}\} \subseteq C_{i3}$ for any i = 1, 2, 3, 4. By the definition of \mathscr{C} , we obtain $C_{13} = C_{23} = C_{33} = C_{43}$ and $\{a_1, a_2, a_3, a_4, a_8, a_9, a_{10}\} \subseteq C_{i3}$ for any i = 1, 2, 3, 4, which is a contradiction to Lemma 3(i).

(b) Put k = 1. Let C_1 and C_2 , respectively, be b_1 - and b_2 -reduced clusters from $\mathscr{C}(a_1)$. Then $C_1 \subseteq \{a_1, a_2, \ldots, a_7\}$, $C_2 = \{a_1\}$. By Lemma 3(i), $|C_1| \leq 6$. Consider the cluster $C_3 \in \mathscr{C}(a_1) \setminus \{C_1, C_2\}$. The equality $C_1 \cup C_2 \cup C_3 = C_1 \cup C_3 = A$ implies that $|C_1| \geq 5$ by Lemma 3(i).

W.l.o.g., assume that $C_1 = \{a_1, a_2, \ldots, a_5\}$. Then $C_3 = \{a_1, a_6, a_7, \ldots, a_{10}\}$ by Lemma 3(i). We obtain $C_1 \setminus C_3 \neq \emptyset$ and $|C_3 \setminus C_1| = 5$, contradicting Lemma 3(ii). Now, w.l.o.g. put $C_1 = \{a_1, a_2, \ldots, a_6\}$. Then $\{a_1, a_7, a_8, a_9, a_{10}\} \subseteq C_3$. By Lemma 3(ii), $|C_1 \setminus C_3| \leq 4$. Therefore, one of the vertices a_2, a_3, \ldots, a_6 , say a_2 , belongs to C_3 . By Lemma 3(i), $C_3 = \{a_1, a_2, a_7, a_8, a_9, a_{10}\}$. Let C_4 be a b_1 reduced cluster from $\mathscr{C}(a_7)$. We get $C_3 \neq C_4$, since $C_3 \notin N(b_1)$. By Lemma 5(i), $|C_4 \cap C_1| \leq 1$. We obtain that $a_7 \in (C_3 \cap C_4) \setminus C_1$ and $|C_1 \setminus (C_3 \cup C_4)| \geq 3$, which is a contradiction to Lemma 3(ii).

(c) Put k = 2. Let C_1 and C_2 , respectively, be b_1 - and b_2 -reduced clusters from $\mathscr{C}(a_1)$. Taking into account the case (b), we can assume that $C_2 = \{a_1, a_2\}$. Then we can proceed analogously to the case (b).

(d) Finally, we assume that k = 3. For any i = 1, 2, 3, denote by C_{i1} and C_{i2} , respectively, b_1 - and b_2 -reduced clusters from $\mathscr{C}(a_i)$. Taking into account the cases (b) and (c), we can assume that $C_{12} = \{a_1, a_2, a_3\}$. Consider the cluster $C_{i3} \in \mathscr{C}(a_i) \setminus \{C_{i1}, C_{i2}\}$. Since $C_{i1}, C_{i2} \subseteq \{a_1, a_2, \ldots, a_7\}$, we have $\{a_8, a_9, a_{10}\} \subseteq C_{i3}$ for any i = 1, 2, 3. By the definition of $\mathscr{C}, C_{13} = C_{23} = C_{33}$ and $\{a_1, a_2, a_3, a_8, a_9, a_{10}\} \subseteq C_{i3}$ for any i = 1, 2, 3. By Lemma 3(i), $C_{i3} = \{a_1, a_2, a_3, a_8, a_9, a_{10}\}$. We obtain that $C_{12} \neq C_{13}$ and $|C_{12} \cap C_{13}| = 3$, which is a contradiction to the definition of \mathscr{C} .

(6) G = G(9, 8, 1).

(a) Assume that there exist exactly two b_1 -reduced clusters $C_1, C_2 \in \mathscr{C}$. Clearly, \mathscr{C} contains a unique b_2 -reduced cluster $C_3 = \{a_1\}$. If $C_1 \cap C_2 = \emptyset$, then $|C_1| = |C_2| = 4$ by Lemma 3(ii). W.l.o.g., assume that $a_1 \in C_1$. Thus, $a_1 \in (C_1 \cap C_3) \setminus C_2$ and $|C_2 \setminus (C_1 \cup C_3)| = 4 > 2$, which is a contradiction to Lemma 3(iii).

Let $C_1 \cap C_2 \neq \emptyset$. It follows from Lemma 5(i) that $|C_1 \cap C_2| = 1$. Then $C_1 \cap C_2 \neq \{a_1\}$ by Lemma 5(ii). Let $C_1 \cap C_2 = \{a_2\}$ and $a_1 \in C_1$. Since $C_1 \notin C_2$ and $C_2 \notin C_1$, we have $|C_1| \leq 5$ and $|C_2| \leq 5$ by Lemma 3(ii). The equality $|C_1 \cup C_2| = 8$ implies $|C_1| \geq 4$ and $|C_2| \geq 4$. We have $a_1 \in (C_1 \cap C_3) \setminus C_2$ and $|C_2 \setminus (C_1 \cup C_3)| \geq 3$, which is a contradiction to Lemma 3(ii).

(b) Now, let C_1, C_2, C_3 and $C_4 = \{a_1\}$, respectively, be three b_1 - and a unique b_2 -reduced clusters in \mathscr{C} . W.l.o.g., assume that $a_1 \in C_1$. By Lemma 5(iii), $C_1 \cap C_i \neq \{a_1\}$ for any i = 2, 3.

Furthermore, we have $|C_1| \ge 5$. Otherwise, $|A \setminus C_1| \ge 5$ and, by the definition of \mathscr{C} , there exists a cluster $C_5 \in \mathscr{C}(a_1) \setminus \{C_1, C_4\}$ such that $(A \setminus C_1) \cup \{a_1\} \subseteq C_5$. By Lemma 3(i), it follows that $|A \setminus C_1| = 5$, i.e., $|C_1| = 4$. We have $C_1 \setminus C_5 \neq \emptyset$ and $|C_5 \setminus C_1| \ge 5$, which is a contradiction to Lemma 3(ii). Therefore, by the same lemma, $C_1 \cap C_2 \neq \emptyset$ and $C_1 \cap C_3 \neq \emptyset$. By Lemmas 5(ii) and 3(ii), we have $|C_1 \setminus C_2| \le 4$ and, consequently, $|C_1| = 5$.

The equality $C_2 \cap C_3 = \emptyset$ holds. Otherwise, by Lemma 5(i) and (iii), we have $(C_2 \cap C_3) \setminus C_1 \neq \emptyset$ and $|C_1 \setminus (C_2 \cup C_3)| = 3$, which is a contradiction to Lemma 3(iii).

Let $C_5 \in \mathscr{C}(a_1) \setminus \{C_1, C_4\}$. Since $A \setminus (C_1 \cup C_4) \subset C_5$, we have $|C_5| \geq 5$. Since $|C_1 \cap C_i| = 1$ for any $i = 2, 3, C_2 \cap C_3 = \emptyset$ and $|(C_2 \cup C_3) \setminus C_1| = 3$, one of the clusters C_2 , C_3 , say C_2 , has size 2. So, we have $(C_2 \cap C_5) \setminus C_1 \neq \emptyset$ and $|C_1 \setminus (C_2 \cup C_5)| \geq 3$ both in the case $|C_5| = 6$ (since $C_2 \subseteq C_5$ by Lemma 3(ii)) and in the case $|C_5| = 5$, which is a contradiction to Lemma 3(iii).

(7) We can come to a contradiction for the graph G = G(9, 8, 2) analogously to the graph G = G(9, 8, 1).

4. Proof of Sufficiency of Theorem 2

Let a threshold graph $H = G(p, q_1, q_2, \ldots, q_k)$ with the bipartition (A, B) not contain any of the graphs $K_{1,4}$, G(12,7), G(11,10), G(10,9,9), G(10,9,7), G(10,9,5), G(10,7,k), $k = 1, 2, \ldots, 7$, G(9,8,2), G(9,8,1) as an induced subgraph. By Theorem 1, we have to show that there exists a (3,2)-covering of H.

W.l.o.g., assume that H is a connected non-complete graph. Therefore, H

has a dominating vertex by the definition of H. Furthermore, $|B| \leq 2$, since H does not contain $K_{1,4}$ as an induced subgraph. Thus, we have $H = G(p, q_1)$ or $H = G(p, q_1, q_2)$.

First, we suppose that $|A| = p \ge 14$. Then $q_1 \le 6$, since H does not contain any of the graphs G(11, 10) and G(12, 7) as an induced subgraph. For any vertex $b \in B$, partition the set N(b) into $n_b \le 3$ pairwise disjoint cliques C_i^b each having size at most 2. Obviously, the list of cliques $(C_i^b \cup \{b\} : b \in B, i = 1, ..., n_b)$ together with the clique A gives a desired (3, 2)-covering of H.

If $|A| \leq 7$, then $q_1 \leq 6$ by the maximality of the clique A. Therefore, a desired (3, 2)-covering of H can be constructed as above.

Now, let $8 \leq |A| \leq 13$. Taking into account the above considerations, we can assume that $q_1 \geq 7$.

Let $H = G(p, q_1)$. Since H does not contain any of the graphs G(12, 7) and G(11, 10) as an induced subgraph, it is isomorphic to one of the graphs G(13, 9), G(12, 9), G(12, 8), G(11, 9), G(11, 8), G(11, 7), G(10, 9), G(10, 8), G(10, 7), G(9, 8), G(9, 7), G(8, 7). Clearly, the set of cliques

$$\mathscr{C} = \{\{a_1, a_2, a_3, a_4, a_5, b_1\}, \{a_1, a_6, a_7, a_8, a_9, b_1\}, \{a_1, a_{10}, a_{11}, a_{12}, a_{13}\}, \\ \{a_2, a_3, a_6, a_7, a_{10}, a_{11}\}, \{a_2, a_3, a_8, a_9, a_{12}, a_{13}\}, \{a_4, a_5, a_6, a_7, a_{12}, a_{13}\}, \\ \{a_4, a_5, a_8, a_9, a_{10}, a_{11}\}\}$$

of the graph G(13,9) is one of its (3,2)-coverings. Each of the graphs G(12,9), G(12,8), G(11,9), G(11,8), G(11,7), G(10,9), G(10,8), G(10,7), G(9,8), G(9,7) and G(8,7) is an induced subgraph of G(13,9). Therefore, a desired (3,2)-covering for each of these graphs can be obtained from the covering \mathscr{C} .

Now, let $H = G(p, q_1, q_2)$. Since H does not contain any of the graphs G(12,7), G(11,10), G(10,9,9), G(10,9,7), G(10,9,5), G(10,7,k), k = 1, 2, ..., 7, G(9,8,2) and G(9,8,1) as an induced subgraph, it is isomorphic to one of the graphs G(11,9,8), G(11,9,6), G(11,9,4), G(10,9,8), G(10,9,6), G(10,9,4), G(10,8,8), G(10,8,7), G(10,8,6), G(10,8,5), G(10,8,4), G(10,8,3), G(9,8,8), G(9,8,7), G(9,8,6), G(9,8,5), G(9,8,4), G(9,8,3), G(9,7,7), G(9,7,6), G(9,7,5), G(9,7,4), G(9,7,3), G(9,7,2), G(9,7,1), G(8,7,7), G(8,7,6), G(8,7,5), G(8,7,4), G(8,7,3), G(8,7,2), G(8,7,1). Some of the desired (3,2)-coverings $\mathscr{C}_1, \mathscr{C}_2, \mathscr{C}_3, \mathscr{C}_4$ for the graphs G(11,9,8), G(11,9,6), G(11,9,4), G(9,7,1), respectively, are given below:

- $\mathscr{C}_1 = \{ \{a_1, a_2, a_3, a_4, a_9, b_1\}, \{a_5, a_6, a_7, a_8, a_9, b_1\}, \{a_1, a_2, a_7, a_8, b_2\}, \\ \{a_3, a_4, a_5, a_6, b_2\}, \{a_1, a_2, a_5, a_6, a_{10}, a_{11}\}, \{a_3, a_4, a_7, a_8, a_{10}, a_{11}\}, \\ \{a_9, a_{10}, a_{11}\}\},$
- $\mathscr{C}_2 = \{\{a_1, a_2, a_7, a_9, b_1\}, \{a_3, a_4, a_7, a_8, b_1\}, \{a_5, a_6, a_8, a_9, b_1\}, \\ \{a_1, a_2, a_3, a_4, a_5, a_6, b_2\}, \{a_5, a_6, a_7, a_{10}, a_{11}\}, \{a_1, a_2, a_8, a_{10}, a_{11}\}, \\ \{a_3, a_4, a_9, a_{10}, a_{11}\}\},$

 $\begin{aligned} \mathscr{C}_3 \ = \ & \{\{a_1, a_2, a_3, a_4, a_9, b_1\}, \{a_5, a_6, a_7, a_8, a_9, b_1\}, \{a_1, a_2, a_7, a_8, b_2\}, \\ & \{a_3, a_4, a_5, a_6\}, \{a_1, a_2, a_5, a_6, a_{10}, a_{11}\}, \{a_3, a_4, a_7, a_8, a_{10}, a_{11}\}, \\ & \{a_9, a_{10}, a_{11}\}\}, \end{aligned}$

$$\mathscr{C}_{4} = \{\{a_{1}, a_{2}, a_{3}, a_{4}, a_{5}, b_{1}\}, \{a_{5}, a_{6}, a_{7}, b_{1}\}, \{a_{1}, b_{2}\}, \{a_{1}, a_{2}, a_{6}, a_{7}, a_{8}, a_{9}\}, \{a_{3}, a_{4}, a_{6}, a_{7}\}, \{a_{3}, a_{4}, a_{8}, a_{9}\}, \{a_{5}, a_{6}, a_{7}\}, \{a_{5}, a_{8}, a_{9}\}\}.$$

Each of the remaining graphs G(10, 9, 8), G(10, 9, 6), G(10, 9, 4), G(10, 8, 8), G(10, 8, 7), G(10, 8, 6), G(10, 8, 5), G(10, 8, 4), G(10, 8, 3), G(9, 8, 8), G(9, 8, 7), G(9, 8, 6), G(9, 8, 5), G(9, 8, 4), G(9, 8, 3), G(9, 7, 7), G(9, 7, 6), G(9, 7, 5), G(9, 7, 4), G(9, 7, 3), G(9, 7, 2), G(8, 7, 7), G(8, 7, 6), G(8, 7, 5), G(8, 7, 4), G(8, 7, 2), G(8, 7, 1) is an induced subgraph for some of the graphs G(11, 9, 8), G(11, 9, 6), G(11, 9, 4), G(9, 7, 1). Therefore, a desired (3, 2)-covering for each of the remaining graphs can be obtained from one of the coverings \mathscr{C}_1 , \mathscr{C}_2 , \mathscr{C}_3 , \mathscr{C}_4 .

5. Recognition Algorithm

The proof of sufficiency of Theorem 2 implies the following linear algorithm for recognizing graphs from L_3^2 in the class of threshold graphs.

Algorithm

Input: a connected threshold graph H with bipartition (A, B), where A is a maximal clique in H.

Output: 1 if $H \in L_3^2$, and 0 otherwise.

1. begin

- 2. **if** $B = \emptyset$, i.e., the graph *H* is complete,
- 3. return 1;
- 4. **if** $|B| \ge 3$
- 5. return 0;
- 6. **if** $\deg(b) \le 6$ for every $b \in B$
- 7. return 1;
- 8. **if** $|A| \ge 14$
- 9. **return** 0;
- 10. **if** H contains some of the graphs G(12,7), G(11,10), G(10,9,9), G(10,9,7), G(10,9,5), G(10,7,k), k = 1, 2, ..., 7, G(9,8,2), G(9,8,1)as an induced subgraph
- 11. **return** 0;
- 12. **return** 1;
- 13. **end**.

The complexity of the algorithm in lines 1–9 is at most O(n), where n = |V(H)|. Since the order of the graph H in line 10 is at most 13, this line takes O(1) time.

So, the total complexity of the recognition algorithm is O(n).

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